

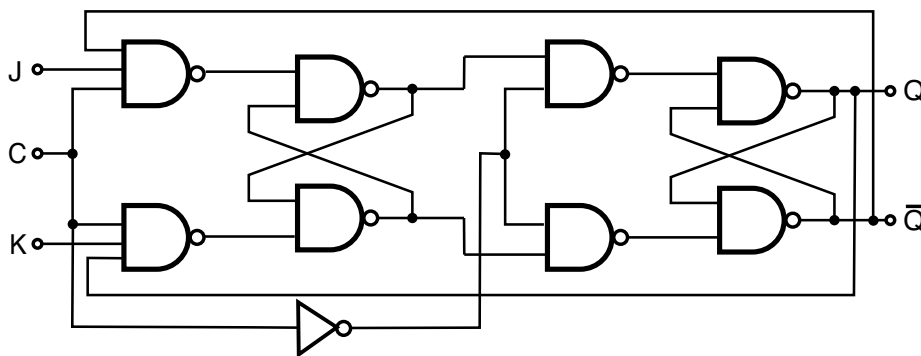
ICS 2021 Problem Sheet #9

Problem 9.1: JK flip-flops

(1+1+1+1 = 4 points)

JK flip-flops, also colloquially known as jump/kill flip-flops, augment the behaviour of SR flip-flops. The letters J and K were presumably picked by Eldred Nelson in a patent application.

The sequential digital circuit shown below shows the design of a JK flip-flop based on two SR NAND latches. Assume the circuit's output is $Q = 0$ and that the inputs are $J = 0$ and $K = 0$, and that the clock input is $C = 0$. (You can make use of the fact that we already know how an SR NAND latch behaves.)

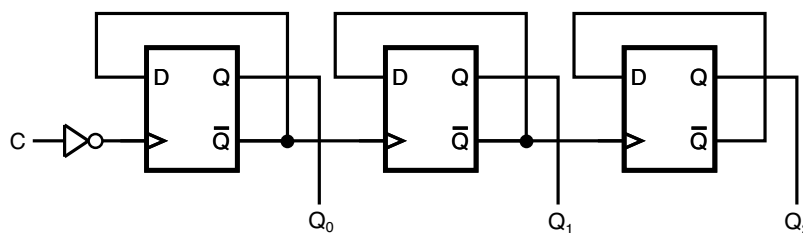


- Suppose J transitions to 1 and C transitions to 1 soon after. Create a copy of the drawing and indicate for each line whether it carries a 0 or a 1.
- Some time later, C transitions back to 0 and soon after J transitions to 0 as well. Create another copy of the drawing and indicate for each line whether it carries a 0 or a 1.
- Some time later, J and K both transition to 1 and C transitions to 1 soon after. Create another copy of the drawing and indicate for each line whether it carries a 0 or a 1.
- Finally, C transitions back to 0 and soon after J and K both transition to 0 as well. Create another copy of the drawing and indicate for each line whether it carries a 0 or a 1.

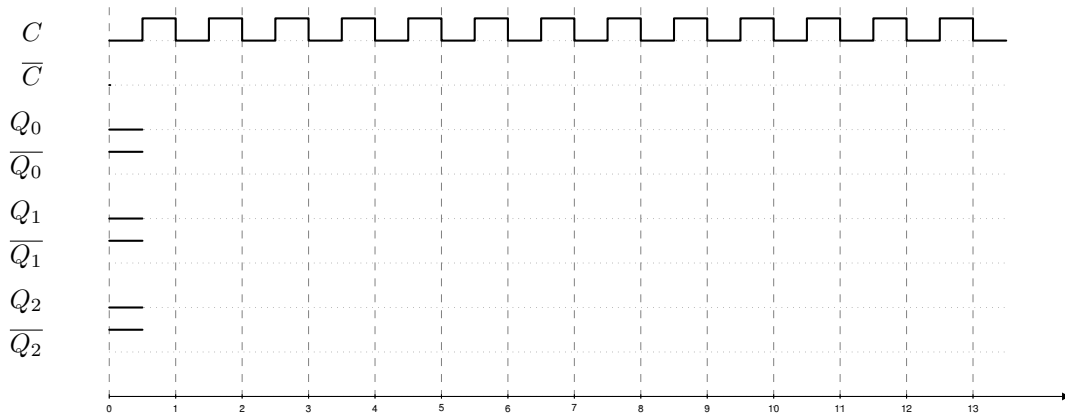
Problem 9.2: ripple counter using d flip-flops

(2+1 = 3 points)

The following circuit shows a 3-bit ripple counter consisting of three positive edge triggered D flip-flops and a negation gate on the clock input C .



- Complete the following timing diagram. Assume that gate delays are very short compared to the speed of the clock signal (i.e., you can ignore the impact of gate delays).



- b) Can you make ripple counters arbitrary “long” or is there a limit on the number of D flip flops that can be chained? Explain.

Problem 9.3: *integer expression simplification (haskell)*

(2+1 = 3 points)

Our goal is to simplify integer expressions that may include constants and variables and which are constructed using a sum and a product operator. For example, we want to write a program that can simplify $(2 \cdot y) \cdot (3 + (2 \cdot 2))$ to $14 \cdot y$. We can represent expressions in Haskell using the following data type:

```

1 data Exp = C Int      -- a constant integer
2           | V String  -- a variable with a name
3           | S Exp Exp -- a sum of two expressions
4           | P Exp Exp -- a product of two expressions
5           deriving (Show, Eq)

```

We use the following rules to simplify expressions:

- S1 Adding two constants a and b yields a constant, which has the value $a + b$, e.g., $3 + 5 = 8$.
- S2 Adding 0 to a variable yields the variable, i.e., $0 + x = x$ and $x + 0 = x$.
- S3 Adding a constant a to a sum consisting of a constant b and a variable yields the sum of the $a + b$ and the variable, e.g., $3 + (5 + y) = 8 + y$.
- P1 Multiplying two constants a and b yields a constant, which as the value $a \cdot b$, e.g., $3 \cdot 5 = 15$.
- P2 Multiplying a variable with 1 yields the variable, i.e., $1 \cdot y = y$ and $y \cdot 1 = y$.
- P3 Multiplying a variable with 0 yields the constant 0, i.e., $0 \cdot y = 0$ and $y \cdot 0 = 0$.
- P4 Multiplying a constant a with a product consisting of a constant b and a variable yields the product of $a \cdot b$ and the variable, e.g., $3 \cdot (2 \cdot y) = 6 \cdot y$.

The usual associativity rules apply. Note that we have left out distributivity rules.

- a) Implement a function `simplify :: Expr -> Expr`, which simplifies expressions that do not contain variables. In other words, `simplify` returns the (constant) value of an expression that does not contain any variables.
- b) Extend the function `simplify` to handle variables as described above.

Submit your Haskell code plus an explanation (in Haskell comments) as a plain text file. Below is a template providing a collection of test cases.

```

1  module Main (main) where
2
3  import Test.HUnit
4
5  data Exp = C Int          -- a constant integer
6           | V String      -- a variable with a name
7           | S Exp Exp     -- a sum of two expressions
8           | P Exp Exp     -- a product of two expressions
9           deriving (Show, Eq)
10
11 simplify :: Exp -> Exp
12 simplify _ = undefined
13
14 tI0 = TestList
15     [ simplify (C 3) ~?= C 3          -- 3 = 3
16       , simplify (V "y") ~?= V "y"   -- y = y
17     ]
18
19 tS1 = TestList
20     [ simplify (S (C 3) (C 5)) ~?= C 8          -- 3 + 5 = 8
21     ]
22
23 tS2 = TestList
24     [ simplify (S (C 0) (V "y")) ~?= V "y"     -- 0 + y = y
25       , simplify (S (V "y") (C 0)) ~?= V "y"   -- y + 0 = y
26     ]
27
28 tS3 = TestList
29     [ simplify (S (S (C 3) (V "y")) (C 5)) ~?= S (C 8) (V "y")   -- (3 + y) + 5 = 8 + y
30       , simplify (S (S (V "y") (C 3)) (C 5)) ~?= S (C 8) (V "y") -- (y + 3) + 5 = 8 + y
31       , simplify (S (C 3) (S (C 5) (V "y"))) ~?= S (C 8) (V "y") -- 3 + (5 + y) = 8 + y
32       , simplify (S (C 3) (S (V "y") (C 5))) ~?= S (C 8) (V "y") -- 3 + (y + 5) = 8 + y
33     ]
34
35 tS4 = TestList
36     [ simplify (S (S (C 3) (C 5)) (C 8)) ~?= C 16          -- (3 + 5) + 8 = 16
37       , simplify (S (C 3) (S (C 5) (C 8))) ~?= C 16       -- 3 + (5 + 8) = 16
38       , simplify (S (C 5) (V "y")) ~?= S (C 5) (V "y")   -- 5 + y = 5 + y
39       , simplify (S (V "y") (C 5)) ~?= S (V "y") (C 5)   -- y + 5 = y + 5
40       , simplify (S (V "x") (V "y")) ~?= S (V "x") (V "y") -- x + y = x + y
41     ]
42
43 tP1 = TestList
44     [ simplify (P (C 3) (C 5)) ~?= C 15          -- 3 * 5 = 15
45     ]
46
47 tP2 = TestList
48     [ simplify (P (C 1) (V "y")) ~?= V "y"     -- 1 * y = y
49       , simplify (P (V "y") (C 1)) ~?= V "y"   -- y * 1 = y
50     ]
51
52 tP3 = TestList
53     [ simplify (P (C 0) (V "y")) ~?= C 0        -- 0 * y = 0
54       , simplify (P (V "y") (C 0)) ~?= C 0     -- y * 0 = 0
55     ]
56
57 tP4 = TestList
58     [ simplify (P (P (C 3) (V "y")) (C 2)) ~?= P (C 6) (V "y") -- (3 * y) * 2 = 6 * y
59       , simplify (P (P (V "y") (C 3)) (C 2)) ~?= P (C 6) (V "y") -- (y * 3) * 2 = 6 * y
60       , simplify (P (C 3) (P (C 2) (V "y"))) ~?= P (C 6) (V "y") -- 3 * (2 * y) = 6 * y
61       , simplify (P (C 3) (P (V "y") (C 2))) ~?= P (C 6) (V "y") -- 3 * (y * 2) = 6 * y
62     ]
63

```

```

64 tP5 = TestList
65   [ simplify (P (P (C 3) (C 5)) (C 8)) ~?= C 120                -- (3 * 5) * 8 = 120
66   , simplify (P (C 3) (P (C 5) (C 8))) ~?= C 120              -- 3 * (5 * 8) = 120
67   , simplify (P (C 5) (V "y")) ~?= P (C 5) (V "y")          -- 5 * y = 5 * y
68   , simplify (P (V "y") (C 5)) ~?= P (V "y") (C 5)          -- y * 5 = y * 5
69   , simplify (P (V "x") (V "y")) ~?= P (V "x") (V "y")      -- x * y = x * y
70   ]
71
72 tM0 = TestList [
73   -- (2 * y) * (3 + (2 * 2)) = 14 * y
74   simplify (P (P (C 2) (V "y")) (S (C 3) (P (C 2) (C 2)))) ~?= P (C 14) (V "y")
75   -- x + (1 + -1) = x
76   , simplify (S (V "x") (S (C 1) (C (-1)))) ~?= V "x"
77   -- (1 + -1) * x = 0
78   , simplify (P (S (C 1) (C (-1))) (V "x")) ~?= C 0
79   -- (2 + -1) * x = x
80   , simplify (P (S (C 2) (C (-1))) (V "x")) ~?= V "x"
81   -- (2 * 2) * (3 + 4) = 28
82   , simplify (P (P (C 2) (C 2)) (S (C 3) (C 4))) ~?= C 28
83   ]
84
85 main = runTestTT $ TestList [tI0, tS1, tS2, tS3, tS4, tP1, tP2, tP3, tP4, tP5, tM0 ]

```